LING 364: Introduction to Formal Semantics

Lecture 25 April 18th

Administrivia

- Homework 5
 - graded and returned

Administrivia

- Today
 - review homework 5
 - also new handout
 - Chapters 7 and 8
 - we'll begin talking about tense

- Exercise 1: Truth Tables and Prolog
- Question A: Using
- ?- implies(P,Q,R1), or(P,Q,R2), \+ R1 = R2.
- for what values of P and Q are P⇒Q and PvQ incompatible?

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?- implies(P,Q,R1), or(P,Q,R2), \+ R1=R2.

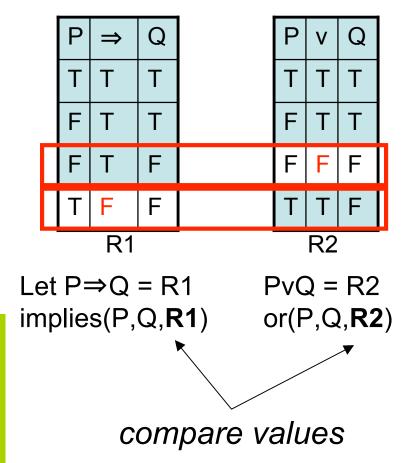
P = true, Q = false,

R1 = false, R2 = true ?;

P = false, Q = false,

R1 = true, R2 = false ?;

no
```



- Exercise 1: Truth Tables and Prolog
- Question B
- Define truth table and/3 in Prolog

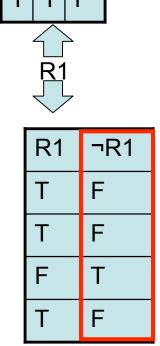
P ^ Q
T T T
F F T
F F
T F F
Result

% and(P,Q,Result) and(true,true,true). and(true,false,false). and(false,true,false). and(false,false,false).

- Exercise 1: Truth Tables and Prolog
- Question C
- Show that
- $\neg (P \lor Q) = \neg P \land \neg Q$
- (De Morgan's Rule)

?- or(P,Q,R1), neg(R1,NR1), neg(P,NP), neg(Q,NQ), and(NP,NQ,R2), \+ NR1=R2.

No



Q

V

Р	¬P			Q	¬Q
Т	F			Т	F
F	Т			Т	F
F	Т			F	Т
Τ	F			F	Т
	¬Р	٨	٦	Q	
	I	F	FF		
	Т	F	F		
	Т	Т	Т		
	F	F	Т		
R2					

- Exercise 1: Truth Tables and Prolog
- Question D
 - Show that
 - $\neg (P \land Q) = \neg P \lor \neg Q$
 - (another side of De Morgan's Rule)
- Question C was for
- $\neg (P \lor Q) = \neg P \land \neg Q$

```
?- and(P,Q,R1),
neg(R1,NR1), neg(P,NP),
neg(Q,NQ), or(NP,NQ,R2),
\+ NR1=R2.
```

- **Exercise 2: Universal**
- **Quantification and Sets**
- **Assume meaning grammar:**

-?- s(M,[every,woman,likes,ice,cream],[]).

```
s(M) \longrightarrow gnp(M), vp(P), \{predicate2(M,P)\}.
  n(woman()) --> [woman].
  vp(M) \longrightarrow v(M), np(X), \{saturate2(M,X)\}.
  v(likes( X, Y)) --> [likes].
  np(ice cream) --> [ice,cream].
  qnp(M) \longrightarrow q(M), n(P), \{predicate1(M,P)\}.
  q((findall(X, P1,L1),findall(Y, P2,L2),subset(L1,L2))) --> [every].
every has semantics \{X: P_1(X)\}\subseteq \{Y: P_2(Y)\}
```

```
saturate1(P,X):- arg(1,P,X).
                                                           saturate2(P,X) :- arg(2,P,X).
                                                           subset([], ).
                                                           subset([X|L1],L2) :- member(X,L2),
                                                               subset(L1,L2).
                                                           member(X,[X|]).
                                                           member(X,[\ |L]) :- member(X,L).
                                                           predicate1((findall(X,P, ), ),P):-
                                                               saturate1(P,X).
                                                           predicate2(( ,(findall(X,P, ), )),P) :-
                                                               saturate1(P,X).
every woman likes ice cream \{X: woman(X)\}\subseteq \{Y: likes(Y, ice_cream)\}
     -M = findall(A, woman(A), B), findall(C, likes(C, ice cream), D), subset(B, D)
```

- Exercise 2: Universal
- Quantification and Sets
- Questions A and B
 - John likes ice cream

Simple way (not using Generalized Quantifiers)

s(P) --> namenp(X), vp(P), {saturate1(P,X)}. namenp(john) --> [john].

note: very different from s(M) --> qnp(M), vp(P), {predicate2(M,P)}.

```
?- s(M,[john,likes,ice,cream],[]).
M = likes(john,ice_cream)
```

```
?- s(M,[john,likes,ice,cream],[]), call(M).
M = likes(john,ice_cream)
```

database

woman(mary).
woman(jill).
likes(john,ice_cream).
likes(mary,ice_cream).
likes(jill,ice_cream).

- Exercise 2: Universal Quantification and Sets
- Question C
 - (names as Generalized Quantifiers)
 - Every woman and John likes ice cream
 - $({X: woman(X)})$ ∪ ${X: john(X)})$ ⊆ ${Y: likes(Y,ice_cream)}$
 - John and every woman likes ice cream

Treat *John* just like *every*:

- Exercise 2: Universal Quantification and Sets
- Question C
 - John and every woman likes ice cream
 - ({X: john(X)} ∪{Y: woman(Y)}) ⊆ {Z: likes(Z,ice_cream)}
 findall P1 findall P2 findall
 union subset

```
Define conjnp:
```

```
s(M) \dashrightarrow \textbf{conjnp}(M), vp(P), \{predicate2(M,P)\}. \\ \textbf{conjnp}((\underbrace{findall(X,\textbf{P1},L1),findall(Y,\textbf{P2},L2),union(L1,L2,L3)),findall(\_,\_,L4),subset}_{(L3,L4))}) \dashrightarrow \\ namenp(\textbf{M1}), [and], qnp(\textbf{M2}), \{predicate1(M1,P1), predicate1(M2,P2), saturate1(P1,X), saturate1(P2,Y)\}.
```

```
I ?- s(M,[john,and,every,woman,likes,ice,cream],[]).
I M = (findall(A,john(A),B), findall(C,woman(C),D),union(B,D,E)),
  findall(F,likes(F,ice_cream),G),subset(E,G)
```

- Exercise 3: Other Generalized Quantifiers
- Question A

```
no: \{X: P_1(X)\} \cap \{Y: P_2(Y)\} = \emptyset
```

No woman likes ice cream

```
\begin{split} &qnp(M) --> q(M), \, n(P), \, \{predicate1(M,P)\}, \\ &q((findall(\_X,\_P1,L1),findall(\_Y,\_P2,L2),subset(L1,L2))) --> [every], \\ &q((findall(\_X,\_P1,L1),findall(\_Y,\_P2,L2),intersect(L1,L2,[]))) --> [no]. \end{split}
```

```
?- s(M,[no,woman,likes,ice,cream],[]).
M = findall(_A,woman(_A),_B),findall(_C,likes(_C,ice_cream),_D),intersect(_B,_D,[])
?- s(M,[no,woman,likes,ice,cream],[]), call(M).
no
```

- Exercise 3: Other Generalized Quantifiers
- Question A

```
some: \{X: P_1(X)\} \cap \{Y: P_2(Y)\} \neq \emptyset
```

- Some women like ice cream (plural agreement)
- *Some woman likes ice cream

```
 \begin{split} &qnp(M) --> q(M), \ n(P), \ \{predicate1(M,P)\}, \\ &q((findall(\_X,\_P1,L1),findall(\_Y,\_P2,L2),subset(L1,L2))) --> [every]. \\ &q((findall(\_X,\_P1,L1),findall(\_Y,\_P2,L2),intersect(L1,L2,L3),\ +L3=[])) --> [some]. \end{split}
```

```
don't have to implement agreement in this exercise, you could just add: n(woman(_)) --> [women]. v(likes(_X,_Y)) --> [like].
```

- Exercise 3: Other Generalized Quantifiers
- Question A

```
some: \{X: P_1(X)\} \cap \{Y: P_2(Y)\} \neq \emptyset
```

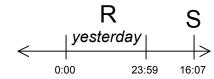
- Some women like ice cream (plural agreement)
- *Some woman likes ice cream

```
q((findall(_X,_P1,L1),findall(_Y,_P2,L2),subset(L1,L2))) --> [every].
q((findall(_X,_P1,L1),findall(_Y,_P2,L2),intersect(L1,L2,L3),\+L3=[])) --> [some].
?- s(M,[some,women,like,ice,cream],[]), call(M).
M =
findall(_A,woman(_A),[mary,jill]),findall(_B,likes(_B,ice_cream),[john,mary,jill]),intersect([mary,jill],[john,mary,jill],[mary,jill]),\+[mary,jill]=[]
```

Chapter 8: Tense, Aspect and Modality

Tense

- Formal tools for dealing with the semantics of tense (Reichenbach):
 - use the notion of an event
 - relate
 - utterance or speech time (S),
 - event time (E) and
 - reference (R) aka topic time (T)



- S,E and T are time intervals:
 - · think of them as time lines
 - · equivalently, infinite sets containing points of time
- examples of relations between intervals:
 - precedence (<), inclusion (⊆)

Past Tense

- Example:
 - (16) Last month, I went for a hike
 - S = utterance time
 - E = time of hike
- What can we infer about event and utterance times?
 - E is within the month previous to the month of S
 - (Note: E was completed last month)
- Tense (went)
 - past tense is appropriate since E < S
- Reference/Topic time?
 - -may not seem directly applicable here
 - $-T = last_month(S)$
 - -think of last_month as a function that given utterance time S
 - -computes a (time) interval
 - -name that interval T

Past Tense

- Example:
 - (16) Last month, I went for a hike
- What can we infer?
 - T = reference or topic time
 - T = last_month(S)
 - E ⊆ T
 - E is a (time) interval, wholly contained within or equal to T
- Tense (went)
 - past tense is appropriate when
 - T < S, E ⊆ T

Past Tense

- Example:
 - (17) Yesterday, Noah had a rash
- What can we infer?
 - T = yesterday(S)
 - "yesterday" is relative to utterance time (S)
 - E = interval in which Noah is in a state of having a rash
 - E may have begun before T
 - E may extend beyond T
 - E may have been wholly contained within T
 - $E \cap T \neq \emptyset$
- Tense (had)
 - appropriate since T < S, E ∩ T $\neq \emptyset$

expression reminiscent of the corresponding expression for the generalized quantifier *some*

Simple Present Tense

In English

```
(18a) Mary runs (simple present) has more of a habitual reading
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does **not** imply

(18b) Mary is running (present progressive)

T = S, run(mary) true @ T

@ T = "at time T"

- i.e. Mary is running right now at utterance time
- (cf. past: T < S)
- However, the simple present works when we're talking about "states"
- Example: (has)
 - (18c) Noah has a rash (simple present)
 - rash(noah) true @ T, T=S
 - i.e. Noah has the property of having a rash right now at utterance time

English simple present tense:

T=S, E has a stative interpretation, E \cap T $\neq \emptyset$

Simple Present Tense

- Some exceptions to the stative interpretation idea
- Historical Present
 - present tense used to describe past events
 - Example:
 - (19a) This guy comes up to me, and he says, "give me your wallet"
 - cf. This guy came up to me, and he said...
- Real-time Reporting
 - describe events concurrent with utterance time
 - Example:
 - (19b) She kicks the ball, and it's a goal!
 - cf. She is kicking the ball